

Comparing Social Network Dynamic Operators

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1. Threshold-driven Social Network Dynamics

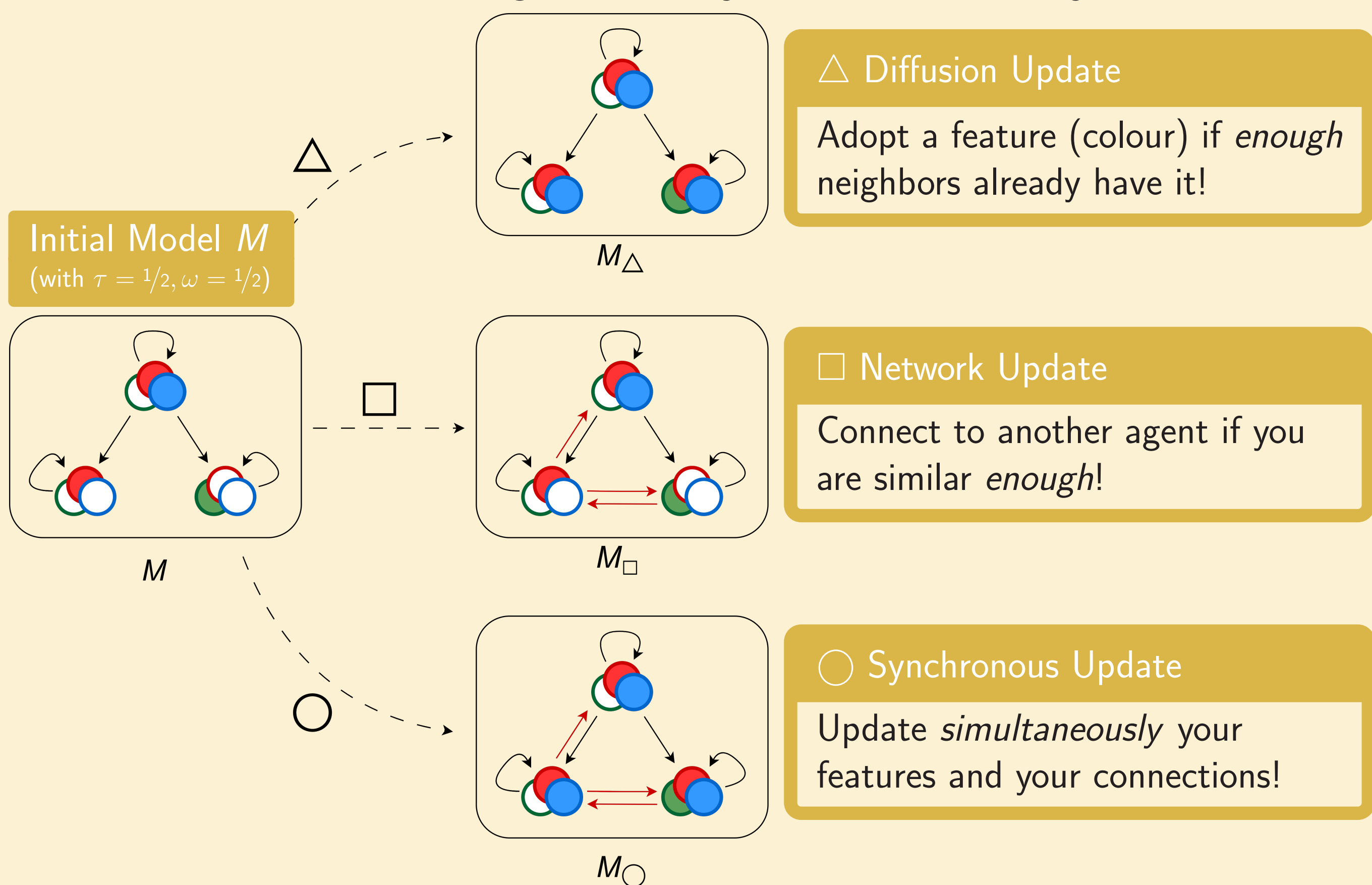
Two main processes can affect agents in a social network:

- **diffusion processes**, opinions spread from one agent to another;
- **similarity network changes**, new links form between similar agents.

Model. Let \mathcal{A} be a non-empty finite set of agents, \mathcal{F} be a non-empty finite set of features. A model M over \mathcal{A} and \mathcal{F} is a tuple $\langle \mathcal{N}, \mathcal{V}, \omega, \tau \rangle$, where:

- $\mathcal{N} \subseteq \mathcal{A} \times \mathcal{A}$ is a social influence relation;
- $\mathcal{V} : \mathcal{A} \rightarrow \mathcal{P}(\mathcal{F})$ assigns to each agent a set of adopted features;
- $\omega \in \mathbb{Q}$ such that $0 \leq \omega \leq 1$ is the *similarity threshold*;
- $\tau \in \mathbb{Q}$ such that $0 < \tau \leq 1$ is the *influenceability threshold*.

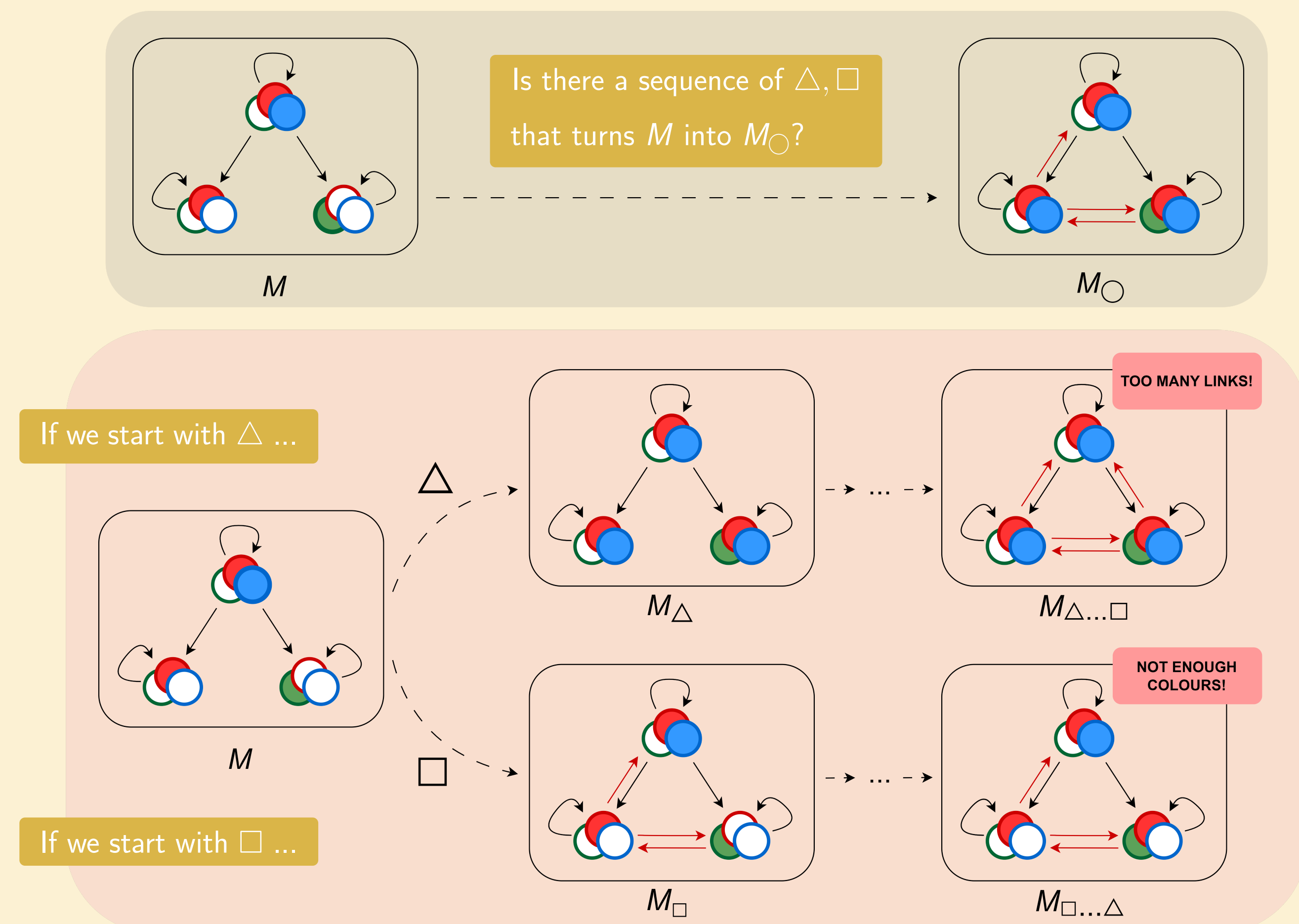
Diffusion and network change can be **synchronous** or **asynchronous**.



3. (Ir)replaceability of Synchronous Operators

QUESTION 1: Is it *always* possible to replace sequences of \circ (synchronous changes) with sequences of \square, Δ (asynchronous changes)?

ANSWER: No! Counterexample:



Sequences of operators. Let $D = \{\circ, \Delta, \square\}$. For $O \subseteq D$, S_O denotes the set of all non-empty finite sequences of operators in O .

Equivalence of sequences. Two sequences $s_1, s_2 \in S_D$ are equivalent on a model M when for all $\varphi \in \mathcal{L}$, $M \models s_1\varphi$ if and only if $M \models s_2\varphi$.

Replaceability (in a model). Let S_1, S_2 be two sets of sequences. S_1 is replaceable with S_2 in a model M , when, for all sequences $s_1 \in S_1$, there exists $s_2 \in S_2$ equivalent to s_1 in M . S_1 is replaceable (tout court) with S_2 when it is replaceable with S_2 in all models.

Theorem (\circ can not *always* be replaced.)

$S_{\{\circ\}}$ is not replaceable with $S_{\{\square, \Delta\}}$.

QUESTION 2: Is it *sometimes* possible to replace \circ ?

ANSWER: Yes! Exactly in these four cases:

- The network is stable.
 $\psi_\Delta := \bigwedge_{a,b \in \mathcal{A}} (N_{ab} \vee \neg \text{sim}_{ab}^\omega)$
- The diffusion process is stable.
 $\psi_\square := \bigwedge_{a \in \mathcal{A}} \bigwedge_{f \in \mathcal{F}} (f_a \vee \neg f_{N(a)}^\tau)$
- A diffusion update preserves similarity between any two agents.
 $\psi_{\Delta \square} := \bigwedge_{a,b \in \mathcal{A}} (\neg N_{ab} \rightarrow (\text{sim}_{ab}^\omega \leftrightarrow \Delta \text{sim}_{ab}^\omega))$
- New links do not prevent conforming to old pressures.
 $\psi_{\square \Delta^n} := \bigwedge_{a \in \mathcal{A}} \bigwedge_{f \in \mathcal{F}} (\neg f_a \rightarrow (f_{N(a)}^\tau \leftrightarrow \bigvee_{0 \leq i \leq n-1} \square \Delta^i f_{N(a)}^\tau))$

Theorem (When can \circ be replaced?)

\circ is replaceable by $S_{\{\Delta, \square\}}$ on a model M iff
 $M \models \psi_\Delta \vee \psi_\square \vee \psi_{\Delta \square} \vee \bigvee_{0 < n < |\mathcal{A}|} \psi_{\square \Delta^n}$.

Theorem (By what can \circ be replaced?)

If $s \in S_{\{\square, \Delta\}}$ is equivalent to \circ on a model M , then s is equivalent to a sequence in the set $\{\square, \Delta, \Delta \square\} \cup \{\square \Delta^n : n > 0\}$ on M .

QUESTION 3: Is it *sometimes* possible to replace \circ^m ?

ANSWER: Yes! At least in this case:

$$M \models \bigwedge_{0 \leq i \leq (m-1)} \circ^i (\psi_\Delta \vee \psi_\square \vee \psi_{\Delta \square} \vee \bigvee_{0 < n < |\mathcal{A}|} \psi_{\square \Delta^n}).$$

2. A Logic of Asynchronous and Synchronous Changes

Syntax \mathcal{L} . The syntax \mathcal{L} is given by

$$\varphi := N_{ab} \mid f_a \mid \neg \varphi \mid \varphi \wedge \psi \mid \Delta \varphi \mid \square \varphi \mid \circ \varphi$$

where $f \in \mathcal{F}$ and $a, b \in \mathcal{A}$.

Conformity pressure. Agent a has network pressure to adopt feature f :

$$f_{N(a)}^\tau := \bigvee_{\{G \subseteq N(a), N \neq \emptyset : \frac{|G|}{|N|} \geq \tau\}} \left(\bigwedge_{b \in N} N_{ba} \wedge \bigwedge_{b \notin N} \neg N_{ba} \wedge \bigwedge_{b \in G} f_b \right).$$

Similarity pressure. Agent a has similarity pressure to connect to b :

$$\text{sim}_{ab}^\omega := \bigvee_{\{E \subseteq \mathcal{F} : \frac{|E|}{|\mathcal{F}|} \geq \omega\}} \bigwedge_{f \in E} (f_a \leftrightarrow f_b).$$

Reduction axioms. Fix $\omega \in [0, 1]$ and $\tau \in (0, 1]$.

$\square N_{ab} \leftrightarrow N_{ab} \vee \text{sim}_{ab}^\omega$	$\Delta N_{ab} \leftrightarrow N_{ab}$	$\circ N_{ab} \leftrightarrow N_{ab} \vee \text{sim}_{ab}^\omega$
$\square f_a \leftrightarrow f_a$	$\Delta f_a \leftrightarrow f_a \vee f_{N(a)}^\tau$	$\circ f_a \leftrightarrow f_a \vee f_{N(a)}^\tau$
$\square(\varphi \wedge \psi) \leftrightarrow \square \varphi \wedge \square \psi$	$\Delta(\varphi \wedge \psi) \leftrightarrow \Delta \varphi \wedge \Delta \psi$	$\circ(\varphi \wedge \psi) \leftrightarrow \circ \varphi \wedge \circ \psi$
$\square \neg \varphi \leftrightarrow \neg \square \varphi$	$\Delta \neg \varphi \leftrightarrow \neg \Delta \varphi$	$\circ \neg \varphi \leftrightarrow \neg \circ \varphi$

Logic $L^{\omega, \tau}$. For fixed ω, τ , the logic $L^{\omega, \tau}$ is some complete axiomatisation and derivation rules of propositional logic, with the above reduction axioms and substitution of provably equivalents.

Theorem (Completeness)

Fix $\omega \in [0, 1]$ and $\tau \in (0, 1]$. For any $\varphi \in \mathcal{L}$:

$$\models_{\mathcal{C}^{\omega, \tau}} \varphi \text{ iff } \vdash_{L^{\omega, \tau}} \varphi.$$

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